

# MICROPROCESSOR & INTERFACING LAB NOTES

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## Microprocessor and Interfacing Lab Notes, First Edition

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The course content was prepared during Spring, 2013.

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18.2.2013

# Lab - 1

## Introduction to debug

-h

AX=0000 BX=0000 CX=0000 DX=E000 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378  
IP=0100  NV  UP  EI  PL  NZ  NA  PO  NC  
1378:0100 0000    ADD [BX+SI], AL  
                  DS:0000 = CD

-?

assemble	A	[address]
compare	C	range address
dump	D	[range]
enter	E	address [list]
fill	F	range list
go	G	[= address] [addresses]
hex	H	value 1 value 2
input	I	port
load	L	[address] [drive] [first sector] [number]
move	M	range address
name	N	[pathname] [arglist]
output	O	port byte

proceed P [= address] [number]  
quit Q  
register R [register]  
search S range list  
trace T [= address] [value]  
unassemble U [range]  
write W [address] [drive] [first sector] [number]  
allocate expanded memory XA [# pages]  
deallocate expanded memory XD [handle]  
map expanded memory pages XM [page] [page] [handle]  
display expanded memory status XS

2. Move 9C to the lower bit of AX register and move 64 to higher bit of DX register. Add AX and DX registers. Note the flag changes and explain in your record.

- R

AX = 0000 BX = 0000 CX = 0000 DX = 0000 SP = FFEE  
BP = 0000 SI = 0000 DI = 0000 DS = 1378 ES = 1378  
SS = 1378 IP = 0100 NV UP EI PL NZ NA PO NC  
1378:0100 0000 ADD [BX+SI], AL DS:0000=C0

- a cs:100

1378:0100 mov ax, 9C  
1378:0103 mov dx, 64  
1378:0105 add ax, dx  
1378:0107

- t = 100 3

AX = 009C BX = 0000 CX = 0000 DX = 0000 SP = FFEE BP = 0000  
SI = 0000 DI = 0000 DS = 1378 ES = 1378 SS = 1378 IP = 1003  
IP = 0103 NV UP EI PL NZ NA PO NC  
1378:0103 B664 MOV DH, 64

AX=009C BX=0000 CX=0000 DX=6400 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378  
IP=0105 NV UP EI PL NZ NA PO NC  
1378: 0105 01D0 ADD AX, DX

AX=649C BX=0000 CX=0000 DX=6400 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378  
IP=0107 NV UP EI PL NZ NA PE NC  
1378: 0107 0000 ADD [BX+SI], AL.  
DS: 0000 = CD

- Seeing the changes in the flag, we find, only the Parity flag changing from PO (odd) to PE (even). This occurs due to the presence of even number of 1's in the final (last) place of data storage i.e. AX.

As shown above, AX = 649C.

Converting to binary,  $649C = 0110\ 0100\ 1001\ 1100$

Seeing the last 8 bits, the number of 1's are even. Hence, the parity.

Move 5 to higher bit of BX register and move 6 to lower bit of CX register. Add both registers and show result in AX register.

- a CS:100

1378:0100 mov bh,5

1378:0102 mov cx,6

1378:0105 add bx,cx

1378:0107 mov ax,bx

1378:0109

- t = 100 4

AX=0000 BX=0500 CX=0000 DX=0000 SP=FFEE BP=0000

SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP=0102 NV UP EI PL NZ NA PO NC

1378:0102 B90600 MOV CX,0006

AX=0000 BX=0500 CX=0006 DX=0000 SP=FFEE BP=0000

SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP=0105 NV UP EI PL NZ NA PO NC

1378:0105 01CB ADD BX,CX

AX=0000 BX=0506 CX=0006 DX=0000 SP=FFEE BP=0000

SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP=0107 NV UP EI PL NZ NA PE NC  
1378:0107 89D8 MOV AX, BX

AX=0506 BX=0506 CX=0006 DX=0000 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP=0109 NV UP EI PL NZ NA PE NC  
1378:0109 0000 ADD [BX+SI], AL  
DS:0506=00

-x-

2. Move 5 to the higher bit of BX register and move 6 to the lower bit of CX register. SUB both registers and show result in AX register. Note flag changes and explain in your record.

-a 100

1378:0100 mov bh, 5

1378:0102 mov cx, 6

1378:0105 sub bx, cx

1378:0107 mov ax, bx

1378:0109

-t = 100 4

AX=0000 BX=0500 CX=0000 DX=0000 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP = 0102 NV UP EI PL NZ NA PO NC  
1378:0102 B90600 MOV CX, 0006

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3/3/2013

AX=0000 BX=0500 CX=0006 DX=0000 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP = 0105 NV UP EI PL NZ NA PO NC  
1378:0105 29CB SUB BX, CX

AX=0000 BX=04FA CX=0006 DX=0000 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP = 0107 NV UP EI PL NZ AC PE NC  
1378:0107 89D8 MOV AX, BX

AX=04FA BX=04FA CX=0006 DX=0000 SP=FFEE BP=0000  
SI=0000 DI=0000 DS=1378 ES=1378 SS=1378 CS=1378

IP = 0109 NV UP EI PL NZ AC PE NC  
1378:0109 0000 ADD [BX+SI], AL DS:04FA=00

Observing the flag changes, we find 2 flags - parity and auxiliary carry to be changing.

Now,  $AX = (04FA)_{16} = (0000\ 0100\ 1111\ 1010)_2$

No. of 1's in AX are 7. And, there is one more 1 for AC. So, total number of 1's are  $7+1=8$ .

Hence, the flag changes from PO (odd parity) to PE (even parity). No. of 1's in last 8 bits are even. Hence, the parity.

25.2.2013

## Lab - 2

### Addressing Modes

Q.1. Write an assembly language program to move 9f2 to AX and 89 to BX respectively, add and store the result in AX.

```
- a cs:300
137F:0300 mov ax, 9f2
137F:0303 mov bx, 89
137F:0306 add ax, bx
137F:0308
- t = 300 3
```

```
AX=09F2 BX=00F9 CX=0012 DX=0000 SP=FFEE BP=0000
SI=0000 DI=0000 DS=137F ES=137F SS=137F CS=137F
IP=0303  NV UP  EI PL  NZ NA  Po  NC
137F:0303  BB8900  MOV  BX,0089
```

```
AX=09F2 BX=0089 CX=0012 DX=0000 SP=FFEE BP=0000
SI=0000 DI=0000 DS=137F ES=137F SS=137F CS=137F
IP=0306  NV UP  EI PL  NZ NA  Po  NC
137F:0306  01D8  ADD  AX, BX
```

AX=0A7B BX=0089 CX=0012 DX=0000 SP=FFEE BP=0000  
 SI=0000 DI=0000 DS=137F ES=137F SS=137F CS=137F  
 IP=0308 NV UP EI PL NZ NA PE NC  
 137F:0308 0000 ADD [BX+SI], AL  
 DS:0089=42

✗

On adding AX and BX, the parity changes from odd (PO) to even (PE).

Seeing the last 8 bits of result stored in AX, 0A7B,

7B  $\equiv$  0111 1011

Total no. of 1's are even. Hence, the change.

Q.2. Write an assembly language program to move 7812 to CX.

- a CS:400

137F:0400 mov CX, 7812

137F:0403

- b = 400

DX=0000

AX=0A7B BX=0089 CX=7812, SP=FFEE BP=0000

SI=0000 DI=0000 DS=137F ES=137F SS=137F

CS=137F IP=0403 NV UP EI PL NZ NA PO CY

137F:0403:0403 0000 ADD [BX+SI], AL

DS:0089=4C

Q.3. Write an assembly language program to move the data stored in address location 4050 to bh register.

```
- e4050
137F:04050 AA:bb
- a cs:300
137F:0300 mov bh, [4050]
137F:0304
- t = 300
```

```
AX=0000 BX=BB00 CX=0000 DX=0000 SP=FFEE BP=0000
SI=0000 DI=0000 DS=137F ES=137F SS=137F CS=137F
IP=0304 NV UP EI PL NZ NA PO CY
137F:0304 0000 ADD [BX+SI], AL
DS:BB00=00
```

\* Different addressing modes used in this lab:

1. Register addressing
2. Immediate addressing
3. Direct addressing

1. Register addressing  
Most 8086 instructions can operate on the 8086's general purpose register set.

By specifying the name of the register as an operand to instruction, you may access the contents of that register.

mov AX, BX: copies value from BX into AX

mov DL, AL: copies value from AL into DL

mov SI, DX: copies value from DX into SI

mov AX, AX: yes, this is legal.

## 2. Immediate addressing

Source operand can be immediate data (8 bit or 16 bit data).

mov AH, 12H

mov CX, 1234H

(note that changes done to 8 bits of the 16 bit register does not affect the other 8 bits of the register)

## 3. Direct addressing

The direct addressing, directly specifies the address as of one of the operand

mov AL, ~~DS~~ [2088H]

loads the AL register with a copy of the byte at memory location.

4.3.2013.

LAB-2 Continued

Q.4. Write an assembly language program to mov 35 to AX and 25 to BX respectively. Subtract and store the result in AX.

```
- a cs:100
13A4:0100 mov ax,35
13A4:0103 mov bx,25
13A4:0106 sub ax,bx
13A4:0100
- t = 100 3
```

```
AX = 0035 BX = 0000 CX = 0000 DX = 0000 .....
IP = 0103 NV UP EI PL NZ NA PO NC
13A4:0103 BB2500 MOV BX,0025
```

```
AX = 0035 BX = 0025 - - - - -
IP = 0106 NV UP EI PL NZ NA PO NC
13A4:0106 29D8 SUB AX,BX
```

```
AX = 0010 BX = 0025 - - - - -
IP = 0100 NV UP EI PL NZ NA PO NC
13A4:0100 0000 ADD [BX+SI],AL
DS:00025 = FF
```

Here, there are no flag changes. The parity (PO) is odd finally as well.

This is because,  $AX = 0035 = 0011\ 0101$ , for last 8 bits, number of 1's is odd.

Similarly, for  $AX = 0010 = 0001\ 0000$ .

So, no. of 1's is again odd.

So, no change in parity.

Final result is got as:  $AX = 0035$

$BX = 0025$

Sub

finally  $AX = 0010$

Q.5. Write an assembly language program to move the data stored in address location 3050 to bh register and add with AL register. Where  $AL = 20$  and memory location  $3050 = 30$ .

- e3050

13A4:3050 00.30

- a cs:200

13A4:0200 mov al, 20

13A4:0202 mov bh, [3050]

13A4:0206 add al, bh

13A4:0208

- t = 200 3

AX = 0200 BX = 0025 - - - - -  
 -- IP = 0202 NV UP EI PL NZ NA PO NC  
 13A4:0202 8A3E5030 MOV BH, [3050]  
 DS:3050 = 30.

AX = 0020 BX = 3025 - - - - -  
 -- IP = 0206 NV UP EI PL NZ NA PO NC  
 13A4:0206 00F8 ADD AL, BH

AX = 0050 BX = 3025 - - - - -  
 -- IP = 0200 NV UP EI PL NZ NA PE NC  
 13A4:0200 0000 ADD [BX+SI], AL  
 DS:3025 = 00

Here, as asked, the value 30 is stored in the address 3050.

Performing the addition operation, we find flag changes from PO to PE.

We had BX = 3025 = 0010 0101

The number of 1's for last 8 bits is even. Hence, the parity -

Q.6. From the 2<sup>nd</sup> program data, write an assembly language program to move the data stored in [BX] to AL register.

- e 3025

13A4: 3025 50.30

- a cs: 200

13A4: 0200 mov al, [bx]

13A4: 0202

- t = 200 1

AX = 0030 BX = 3025 - - - - -

- - - - IP = 0202 NV UP EI PL NZ NA P~~E~~ NC

13A4: 0202 8A3E5030 MOV BH, [3050]

AX = 0030 BX = 3025 - - - - -

- - - - IP = 0206 NV UP EI PL NZ NA P~~E~~ NC

13A4: 0206 00F8 MOV AL, [BX]

DS: 3050 = 30

In this question, the value stored in the address 3025 (= 30) gets stored or moved to AL. So, parity <sup>doesn't</sup> change from even to ~~odd~~. we had ~~3025~~ BX = 3025 = 0010 0101. Seeing no. of 1's, it's even. Now, AX = 0030 = 0011 0000. So, no. of 1's is again even.

\* Addressing mode used in Q.6 :

• Register Indirect Addressing

The 80x86 CPUs let you access memory indirectly through a register using the register indirect addressing modes. Register indirect addressing allows data to be addressed at any memory location through an offset address held in any following register:

BX, BP, SI, DI.

mov AL, [BX]

mov AL, [BP]

mov AL, [SI]

mov AL, [DI]

The [BX], [SI] and [DI] modes, use data segment by default. The [BP] addressing mode uses the stack segment (SS) by default. Segment Overriding can also be done.

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## Sample Program 1: output

AX=1000 ... DS=1000 ... ES=1000 IP=001E  
NV UP EI NG NZ NA PO NC

The output shows how the mov instruction transfers the data between registers.

The value of the memory, data 2 i.e., 1000, goes to AX, which is moved to ES. This is confirmed from the got output.

The memory locations - data 1, data w, data b, data 3, data 4 were not used in program output.

18.3.2013

# Lab-3

## MASM: An Overview

Example: Sample Program 1.

- MODEL SMALL

- ~~386~~ data

data 1 db 23h

data 2 dw 01000h

data w dw 50 dup(?)

data b db 10 dup(?)

data 3 db 'a'

data 4 db 11000100b

- CODE

- STARTUP

mov ax, data 2

mov ds, ax

mov es, ax

- exit

end

## Sample Program 2:

Initial:

AX=0000 BX=FF90 CX=0000 DX=12A3 SP=FF90 BP=0000  
SI=0000 DI=0000 DS=12A3 ES=1286 SS=12A3 CS=129C  
IP=0016 FL=63086 NV UP DI NG NZ NA PE NC

Instruction number:

Output change:

1)

AX=0000 IP=0091

↳ hexadecimal value of 0012

2)

CX=0022 IP=001B

↳ 0022: hexadecimal of 34

3)

AX=0022 IP=001D

4)

AX=0004 IP=0020

5)

BX=9999 IP=0024

6)

CX=270F IP=0028

7)

IP=002C

8)

DX=1299 IP=0030

9)

CX=278B IP=0034

10)

BX=1E8B IP=0038

11)

IP=003B

12)

IP=003F

13)

~~BX=000F IP=0042~~

14)

~~AX=0001 IP=0044~~

15)

~~IP=0047~~

## Example : Sample Program 2

- model small
- data

data1 db 23

data2 dw 9999h

data3 dw 9999

array dw 01, 02, 03, 04, 05, 06, 07, 08

- code

- startup

1) mov ah, 12

2) mov cl, 34

3) mov al, cl

4) mov ax, 04

5) mov bx, data2

6) mov cx, data3

7) mov data1, bl

8) mov dl, data1

9) mov cl, ds:[0520]

10) mov bx, ds:[0520]

11) mov ds:[0520], al

12) mov ds:[0520], cx

~~13) mov bx, offset array~~

~~14) mov ds:[0520], al~~

~~15) mov ds:[0520], cx~~

Instruction number :

Output Change :

16)	BX = 000F	IP = 0042
17)	AX = 0001	IP = 0044
18)		IP = 0047
19)	CX = 0099	IP = 004A
20)		IP = 004C
21)	BX = 1000	IP = 004F
22)	SI = 3000	IP = 0052
23)		IP = 0056
24)	AX = 0099	IP = 005A
25)	AX = 0200	IP = 005E
26)	AX = 0001	IP = 0062
27)		IP = 0065
28)	BX = 000F	IP = 0068
29)	DX = 0001	IP = 006A
30)	AX = 0099	IP = 006C
31)	AX = 0300	IP = 006F
32)	BX = 0099	IP = 0072
33)	DX = 00A8	IP = 0075

16) mov bx, offset array  
17) mov al, [bx]  
18) mov [bp], dl  
19) mov cx, [bp]  
20) mov ax, [bx]  
21) mov bx, 1000H  
22) mov si, 3000H  
23) mov [bx+1000h], cx  
24) mov ax, [si-1000h]  
25) mov ax, array[di+1]  
26) mov ax, array[di]  
27) mov di, 0  
28) mov bx, offset array  
29) mov dx, [bx+di]  
30) mov ax, [bp+di]  
31) mov ax, [bx+di+3h]  
32) mov bx, 3  
33) mov dx, array[bx+di]  
• exit  
end

Q.1) (a) Output

AX = 000A BX = 0002 CX = 0003 DX = 0004 SP = 0000  
BP = 0000 SI = 0000 DI = 0000 DS = 128C ES = 128C  
SS = 129C CS = 129C IP = 0012 FL = 3206  
NV UP EI PL NZ NA PE NC.

The sum of AX, BX, CX, DX registers' value gets stored in AX.

(b) Output

AX = 0018 BX = 0005 CX = 0007 DX = 0009 SP = 0000  
BP = 0000 SI = 0000 DI = 0000 DS = 128C ES = 128C  
SS = 129C CS = 129C IP = 010E FL = 3216  
NV UP EI PL NZ AC ~~IF~~ PE NC

The sum of the hexadecimal values stored in AX, BX, CX and DX register gets added to AX as:

$$AX = AX + BX = 03 + 05 = 08$$

$$AX = AX + CX = 08 + 07 = 0F$$

$$AX = AX + DX = 0F + 09 = 18$$

Here, for bytes of data final value of AX, having 1 will be an auxiliary carry. Its corresponding flag is also high.

25.3.2013

# Lab - 3

continued .

Q.1. (a) WAP to add the values stored in the 4 registers AX, BX, CX, DX .

(b) Perform the same operation using bytes of data using registers AL, BL, CL, DL .

(a) . model tiny

. code

. startup

mov ax, 1

mov bx, 2

mov cx, 3

mov dx, 4

add ax, bx

add ax, cx

add ax, dx

. exit

end

(b) . model tiny

. code

. startup

mov bl, 5

mov cl, 7

mov al, 3

mov dl, 9

add ax, bx

add ax, cx

add ax, dx

. exit

end

Q.2) (a) Output

AX = 0017  
DS = 129E ES = 128C SS = 129E CS = 129C IP = 001E  
FL = 3216 NV UP EI PL NZ AC PE NC

Here, AL being a register of 8 bits, will show auxiliary carry as high when there is a carry after 1st byte. So, as per output,  $10 + 13 = (17)_{16}$  which gives AC.

Also,  $(17)_{16} = (10001)_2$ . So, no. of 1's are even.  
Hence, PE flag is high.

(b) Output

AX = 08FC  
DS = 129E ES = 128C SS = 129E CS = 129C IP = 001E  
FL = 3206 NV UP EI PL NZ NA PE NC

AX = 08FC is got by adding hexadecimal values of 1000 and 1300.

$$\begin{array}{r} (1000)_{10} = (03E8)_{16} \\ (1300)_{10} = (0514)_{16} \\ \hline \text{and} \quad \begin{array}{r} 03E8 \\ + 0514 \\ \hline 08FC \end{array} \end{array}$$

Also,  $(08FC)_{16} = (100011111100)_2$ . Seeing the last 8 bits, no. of 1's are 6 (even). So, PE flag is high.

Q.2. WAP to add the contents of 2 memory locations num1, num2.  
Consider the numbers to be (a) bytes (b) words

(a). model small

• data

num1 db 10

num2 db 13

• code

• startup

mov al, num1

~~mov~~

add al, num2

• exit

end.

(b). model small

• data

num1 dw 1000

num2 dw 1300

• code

• startup

mov ax, num1

add ax, num2

• exit

end

Q 3) (a) Output.

AX = 000A . . . . DS = 129E ES = 128C SS = 129E CS = 129C  
IP = 0029 FL = 3206 NV UP EI PL NZ NA PE NC

Here, AX holds the sum of data 1, data 2, data 3 and data 4. ( $= 1 + 2 + 3 + 4 = 0A$ )

This sum is stored in the memory location, SUM. To access the location of sum, from the output window, note the value of DS and address associated with the last instruction where value of AL is going (i.e. BYTE PTR [0012], AL). Here, DS = 129E. So, in the memory tab, editing the location as, 129E:0012 gives location of SUM. The value of sum gets shown after the address as:  
129E:0012 00 0A . . . . .

(b) Output.

AX = 2710 . . . . DS = 129E ES = 128C SS = 129E CS = 129C  
IP = 0029 FL = 3202 NV UP EI PL NZ NA PO NC

The last instruction on output window:

MOV WORD PTR [0016], AX.

So, going to location 129E:0016, we get output of SUM as: 129E:0016 27 10 . . . . .

Q. 3) WAP to add the contents of 4 memory locations data 1, data 2, data 3, data 4.  
Consider the numbers to be (a) Bytes (b) words

(a) • model small

• data

data 1 db 1

data 2 db 2

data 3 db 3

data 4 db 4

sum db 0

• code

• startup

mov al, data 1

add al, data 2

add al, data 3

add al, data 4

mov sum, al

• exit

end

(b) • model small

• data

data 1 dw 1000

data 2 dw 2000

data 3 dw 3000

data 4 dw 4000

sum dw 0

• code

• startup

mov ax, data 1

~~mov~~ add ax, data 2

add ax, data 3

add ax, data 4

mov sum, ax

• exit

end

Q. 4) (a) Output, finally:

AX = 0028 BX = 0013 CX = 0000... DS = 129E ES = 128C

SS = 129E CS = 129C IP = 0027 FL = 3202

NV UP EI PL NZ NA PO NC

Seeing DS (= 129E) and the offset address of data1 (from the instruction mov BX, 000C), we locate the memory location values as:

129E:000C 05 0A 0F 14 19 1E 23

Each of the array elements are the hexadecimal values of multiples of 5.

(b) Output, finally:

The output for a word data remains same as that of a byte, with the difference of the way, the data is added to the array, data1 (2\*2 bytes = 16 bits of data updated everytime with every operation).

Hence, for the last instruction, data will be stored as

129E:000C ~~05~~ 05 0A 0F 14 19 1E 23 00

23 00: It is the last 16 bit data corresponding to  $(0023)_6 = (35)_{10}$ , which is what is expected.

1.4.2013

# Lab - 3

- continued 2.

Q. 1. WAP to create an array of size 7 and load this array with data elements which are multiple of 5. Consider (a) bytes (b) words.

(a) • model small  
• data  
data1 db 7 dup(?)  
• code  
• startup  
mov ax, 5h  
mov cx, 7  
mov bx, offset data1  
l1:  
mov [bx], ax  
add ax, 5h  
inc bx  
loop l1  
• exit  
end

(b) • model small  
• data  
data1 dw 7 dup(?)  
• code  
• startup  
mov ax, 5h  
mov cx, 7  
mov bx, offset data1  
l1:  
mov [bx], ax  
add ax, 5h  
inc bx  
loop l1  
• exit  
end

## Q.2) (a) Output

AX=0009 BX=0010 CX=0000 DX=129E .... DS=129E CS=129C  
IP=0025 FL=3212 NV UP EI PL NZ AC PO NC

AX: holds final result of alternate elements addition  
 $= 1+3+5 = 0009$ .

BX: holds final address of last element of data 1.  
Address of memory location is got as:

DS: Offset address of data 1 = 129E:000A.

Memory location output: 129E:000A 01 02 03 04 05 ....

CX: A counter which runs 3 times, finally coming to 0.

## (b) Output

AX=0009 BX=0018 CX=0000 .... DS=129E CS=129C  
IP=0028 FL=3206 NV UP EI PL NZ NA PE NC

In the output for word data input, the function/role and/or the value stored in the registers AX, BX and CX remains the same. The difference comes in the way, the word data is present <sup>in memory location</sup> and, its used in the program:

129E:000C 01 00 02 00 03 00 04 00 05 00 ...

The data has been stored as 0001  $\rightarrow$  01 00 .

000C is the offset address of data 1 giving rise to show of memory location contents of data 1.

Q.2) WAP to create an array of size 5 and find the sum of alternate elements of this array. Consider (a) Bytes (b) Word.

(a) • model small  
• data  
data1 db 1,2,3,4,5  
• code  
• startup  
mov bx, offset data1  
mov ax, 0H  
mov cx, 3  
L1:  
add al, [bx]  
inc bx  
inc bx  
loop L1  
• exit  
end

(b) • model small  
• data  
data1 dw 1,2,3,4,5  
• code  
• startup  
mov bx, offset data1  
mov ax, 0H  
mov cx, 3  
L1:  
add ax, [bx]  
inc bx  
inc bx  
loop L1  
• exit  
end

Q.1: Output:

AX = 0105 BX = ~~FF~~FF02 DS = 129E ES = 128C  
CS = 129C IP = 0027 NV UP EI NG NZ NA ~~A~~  
PO NC

BL gets stored with 02 as shown in output.  
The quotient of division gets stored in AL and remainder gets stored in AH.  
hence, AX has the value of 0105 with 01 as remainder and 05 as quotient.

The values moved to ANSQ and ANSR memory locations can be seen by seeing the DS: offset address.

From code view,

for ANSQ : 129E:000F  
ANSR : 129E:0010

15.4.2013

## Lab - 3.

- continued 3.

Q.1. WAP to divide unsigned byte contents of memory location NUMB by NUMB1. Store the quotient and remainder in locations ANSQ and ANSR respectively.

- model small

- data

numb dw 11

numb1 db 2

ansq db 0

ansr db 0

- code

- startup

mov ax, numb

mov bl, numb1

div bl

mov ansq, al

mov ansr, ah

- exit

end.

Q.2 : Output :

$$(-100)_{10} = (FF9C)_{16}$$

$$-100/9 = (-11.11)_{10} = (FFF5)_{16}$$

AX=FFF5 BX=FF09 ... DS=129E ES=128C -- CS=129C  
IP=001E NV UP EI NG NZ NA PO NC.

Clearly, AX has the final result FFF5.  
Value FF9C, stored in DI can be got from the  
address 129E:0002.

Q.3 : Output :

$$1+2+3+4+5+6+7+8+9+10 = (55)_{10} = (37)_{16}$$

AX=0037 BX=001A CX=0000 -- -- DS=129E  
CS=129C IP=0022 NV UP EI PL NZ NA PO NC.

The final value of sum gets stored in AL.  
Value of count decreases from 20 to 0 at the  
end.

Q.2. WAP to divide  $-100/9$ .

- model small

- data

d1 dw -100

d2 db 9

- code

- startup

mov ax, d1

mov bl, 9

idiv bl

- exit

end

Q.3. WAP to find the sum of an array of size 10 words.

- model small

- data

d1 dw 1, 2, 3, 4, 5, 6, 7, 8, 9, 10

- code

- startup

mov cx, 10

mov bx, offset d1

l1:

add al, [bx]

inc bx

loop l1

- exit

end

Q. 4 : Output :

AX = 0505 BX = 0020 CX = 0000 DX = 120A DS = 129E  
CS = 129C IP = 0028 NV UP EI PL NZ ~~AC~~ PO NC.

The sum of elements of array =  
 $1 + 2 + 3 + 4 + 5 + 6 + 7 + 8 + 9 + 10 = (55)_{10}$

To find average,  $\frac{\text{Sum}}{\text{no. of elements}} = \frac{55}{10}$   
 $= 5.5$   
or 5 quotient  
5 remainder

AH : stores remainder }  $\Rightarrow$  AX = 0505, as shown.  
AL : stores quotient }  
The count decreases from 20 to 0.

The auxiliary carry bit became high during the operation  $\because$  addition was going on in 8 bits (AL).

Q.4. WAP to find average of 10 elements stored in an array.

- model small

- data

dl dw 1, 2, 3, 4, 5, 6, 7, 8, 9, 10

- code

- startup

```
mov cx, 20
```

```
mov bx, offset dl
```

```
mov ax, 0
```

l1:

```
add al, [bx]
```

```
inc bx
```

```
loop l1
```

```
mov dl, 10
```

```
div dl
```

- exit

```
end
```

Output : Q.1) (a) Multiply 5 by 17

AX = 0055 BX = 0005 . . . . DS = 1299 CS = 1298  
IP = 001A NV UP EI PL NZ NA PE NC

$$17 \times 5 = (85)_{10} = (55)_{16}$$

It is stored in AX, as shown.

• 386 is required while using SHIFT instruction more than once.

(b) Multiply 5 by 30

AX = 0096 BX = 001E . . . DS = 1288 CS = 1298  
IP = 010A FL = 3216 NV UP EI PL NZ AC PE NC

$$30 \times 5 = (150)_{10} = (96)_{16}$$

It is stored in AX.

While multiplication, there is a carry that comes after 8 bits. So, AC bit becomes high.

The logic is: 30 is stored in AX.  
Then, with AX left shifted by 2, we get  $30 \times 4$ .

Finally,  $(30 \times 4) + 30$  happens.

So,  $30 \times 5$  is the result.

22.4.2013

# Lab-3

- continued 4.

Q.1) Write an assembly language program to multiply two numbers without using MUL instruction

(Multiply 5 by 17 and 30)

(a)

- model small
- code
- 386
- startup

```
mov ax, 5
mov bx, ax
shl ax, 4
add ax, bx
```

• exit  
end

(b)

- model tiny
- code
- 386
- startup

```
mov ax, 30
mov bx, ax
shl ax, 2
add ax, bx
```

• exit  
end

Output : Q.2)

AX = 0083 BX = 006C . . . . DS = 1288 CS = 1298  
IP = 0108 NV UP EI PL NZ NA PO NC

$$EF - 6C = 83.$$

This result is stored in AX, as shown.

Output : Q.3)

AX = 0070 BX = 64E0 DS = 129A CS = 1298 IP = 0025  
FL = 3212 NV UP EI PL NZ AC PO NC

$$100 + 12 = (112)_{10} = (70)_{16}.$$

This value is stored in AX.

To check the value stored in 3050 location,  
go to :

129A : 0BEA ,

which is

DS : offset address .

Q.2) Write an assembly language program to  
mov AX EF and mov BX 6C.  
SUB BX from AX and store result in AX  
register.

- model tiny
- code
- startup

```
mov ax, 0EFh  
mov bx, 06Ch  
sub ax, bx
```

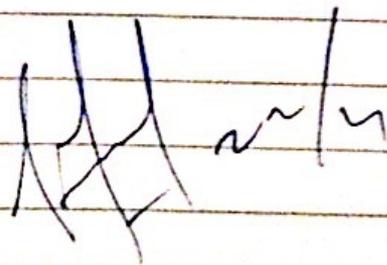
```
• exit  
end
```

Q.3) WAP to move data stored in address location 3050  
to bh register and add with AL register.  
([3050] store 100 and AL store 12)

- model small
- code
- startup

```
mov ds, [3050], 100  
mov al, 12  
mov bh, ds: [3050]  
add al, bh
```

```
• exit  
end.
```



Output : Q.1)

AX=0090 BX=000C SP=0000 IP=0103 DS=128C  
NV UP EI PL NZ NA PE NC

Final result ~~s~~ result on multiplication should be  
 $12 \times 12 = (144)_{10} = (90)_{16}$

That is stored in AX, as shown.

1.5.2013

## Lab-3

- continued 5

Q.1) WAP to perform 16 bit multiplication using procedure.

- model tiny
- code
- startup
- call mult
- exit

```
mult proc near
```

```
mov ax, 12
```

```
mov bx, 12
```

```
mul bx
```

```
ret
```

```
mult endp
```

```
end
```

Output : Q2)

AX = 2D BX = 0005 CX = 0003 DX = 0000 SP = FFEO  
DS = 129E IP = 001A  
FL = 3206 NV UP EI PL NZ NA PE NC

$$BX = BX + CX = 2 + 3 = 5$$

$$AX = BX * (AX - CX)$$

$$= BX * (12 - 3)$$

$$= 5 * 9$$

$$\Rightarrow AX = (45)_{10} = (2D)_{16}$$

This gets stored in AX, as shown.

Q.2) WAP to perform ADD, SUB and MUL using procedure.

- model small
- code
- startup
- call logic
- exit

logic proc near

mov ax, 12

mov bx, 2

mov cx, 3

add bx, cx

sub ax, cx

mul bx

ret

logic endp

end

Q.3) Output :

AX = 0037 BX = FFED CX = 0000 SI = 0012 DS = 13D0  
IP = 0024 NV UP EI PL NZ NA PE NC

$$\begin{aligned} \text{The sum of contents of table} &= 1+2+3+4+5+6+ \\ &\quad 7+8+9+10 \\ &= (55)_{10} \\ &= (37)_{16} \end{aligned}$$

It gets stored in AX, as shown.  
The counter register, CX gets decremented from  
 $(10)_{10} [A]_{16}$  to 0.

3) ~~WAP~~ WAP to add 10 bytes of memory array called TABLE using MACRO.

• model small

• data

table db 01, 02, 03, 04, 05, 06, 07, 08, 09, 10

• code

sums macro NUM, count

local, addrn

mov SI, offset num

mov cx, count

mov AL, 0

addrn:

add AL, [SI]

inc SI

loop addrn

endm

• startup

sums table, 10

• exit

end

Q.1) Output

Procedure

AX = 000A BX = 0019 SP = FFEO IP = 001A FL = 3216  
NV UP EI PL NZ AC PE NC

$$AX = (35)_{10} = (23)_{16}$$

$$BX = (25)_{10} = (19)_{16} \quad : \text{ gets stored in BX}$$

Now,

$$AX - BX = 35 - 25 = (10)_{10} = (000A)_{16}$$

000A is stored in AX, as shown.

Macro

AX = 000A BX = 0019 IP = 001F FL = 3216  
NV UP EI PL NZ AC PE NC ✓

The format/syntax is different as of procedure, but output ~~is~~ is same.

6.5.2013

# Lab - 3

- continued 6

Q.1. Write a procedure and macro for the following -  
to mov 35 to AX and 25 to BX respectively  
subtract and store result in AX.

## Procedure

- model small
- code
- startup
- call move
- exit

```
move proc near
mov ax, 35
mov bx, 25
sub ax, bx
ret
move endp
end
```

## Macro

- model small
- code
- sums macro a, b
- mov ax, a
- mov bx, b
- sub ax, bx
- end m
- startup

```
sums 35, 25
• exit
end
```

## Q.2) Output

### Procedure

AX = 0032 BX = 1EE0 SP = FF00 IP = 001A DS = 129E  
FL = 3212 NV UP EI PL NZ AC PO NC

Value of  $30 + 20 = (50)_{10} = (32)_{16}$ .

This gets stored in AX.  
BH has value of  $(1E)_{16} = (30)_{10}$ .

Also, to access the memory location of 3050,  
we go to 129E : 0BEA.

### Macro

AX = 0032 BX = 1EE0 DS = 129E IP = 0025 FL = 3212  
NV UP EI PL NZ AC PO NC

The same procedure is applied here also, with  
the registers having same values.  
Only the sequence has changed.

Q2. Write a procedure and macro to move the data stored in address location to bh register and add with AL register. (where AL=20 & memory location 3050=30)

### Procedure

- model small
- code
- startup

call logic

- exit

logic proc near

```
mov al, 20
mov ds:[3050], 30
mov bh, ds:[3050]
add al, bh
ret
logic end p
end
```

### Macro

- model small
  - code
- sums macro b b  
add al, b

endm

```
• startup
mov al, 20
mov ds:[3050], 30
mov bh, ds:[3050]
sums bh
```

- exit
- end

~~6/9/2013~~

Q1: Output

AX = 0161 ... SP = 0000 ... IP = 0103 FL = 3202  
NV UP EI PL NZ NA PO NC

Input given: a

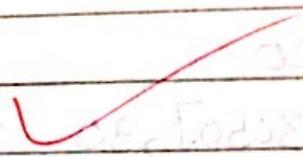
So,  $a = (0097)_{ASCII} = (0061)_{16}$

First 01 goes to AH

next 61 goes to AL (value from INT21h)

So, AX = 0161, as shown.

The output is displayed on the screen  
(<sup>00</sup> with echo)



13.5.13

## Lab - 3

- continued 7

Q.1: WAP to enter character using DOS interrupts with echo.

Ans:

- model tiny
- code
- startup
- call sums
- exit

sums proc near

mov ah, 1

int 21h

or al, al

jnz l1

int 21h

stc

l1: ret

sums endp

end

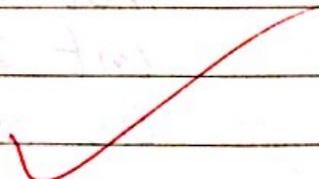
Q.2) Output :

AX = 0761 ... SP = 0000 ... IP = 0103 FL = 3202  
NV UP EI PL NZ NA PO NC

We give 07, stored in AH.  
& value entered from keyboard was the character  $\rightarrow a = (97)_{\text{ASCII}} = (61)_{16}$

This gets stored in AL.  
So, AX = AH AL = 0761, as shown.

the value entered (a) isn't shown on the screen.



Q.2) WAP to enter a character using DOS language interrupts without echo.

- model tiny
- code
- startup

• exit      call sums

sums proc near

MOV AH, 07

MOV DL, 0FFH

INT 21h

JE SUMS

OR AL, AL

JNZ L1

INT 21h

STC

L1: SET

sums endp

end

Q.3) Output :

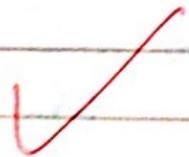
Output is got by directly typing the name of program (instead of going into Code View)

```
masm p1.asm  
link p1  
p1
```

; p1: name of program.

Output : ~~here~~ where is it?

Note: \$ sign after program is necessary



Q.3) WAP to input line from standard input.

- model small

- data

```
mes db 13, 10, 'which year is it? $!'
```

```
mes1 db 10, 10
```

- code

- startup

```
mov ah, 9  
mov dx, offset mes  
int 21h
```

```
mov ah, 0ah  
mov dx, offset mes1  
int 21h
```

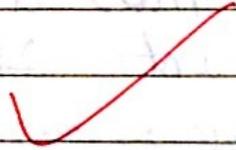
- exit

```
end
```

Q.4) Output :-

On typing the name of the program,  
the screen gets cleared.

The size of clearance is  $25 \times 80$ .



Q.4) WAP to clear monitor using BIOS interrupt

- model small
  - code
  - startup
- dcl

```
mov ax, 0b800h
```

```
mov es, ax
```

```
mov di, 0h ✓
```

```
mov cx, 25*80
```

```
mov ax, 0720h
```

```
rep stosw
```

• exit

end

Q.5) Output :

Output is got by directly writing name of program. (after link)

p5

Output: Akshansh  
          ↓ (enter key)  
Welcome Akshansh ✓

(Akshansh & Welcome Akshansh are typed on the screen.

They become visible side by side)  
Note :- Akshansh gets over written by  
Welcome Akshansh (∵ screen spacing  
25\*80 isn't given)

Q.5) WAP to enter a name and then print welcome along with the name.

- model small

- data

```
buf1 db 257 dup(?)
```

```
buf2 db 257 dup(?)
```

- code

- startup

```
mov buf1, 255
```

```
mov dx, offset buf1
```

```
call line
```

```
mov buf2, 255
```

```
mov dx, offset buf2
```

```
call line
```

- exit

```
line proc near
```

```
mov ah, 0ah
```

```
int 21h
```

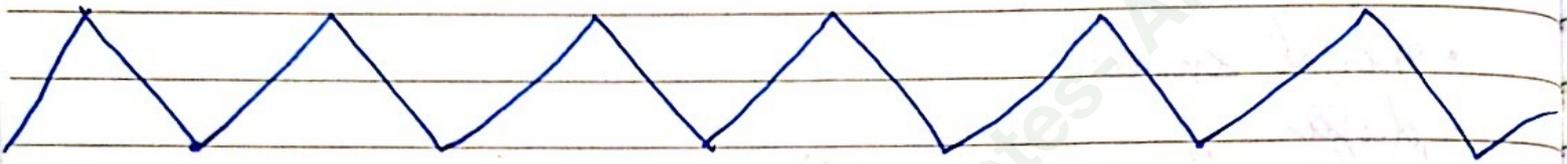
```
ret
```

```
line endp
```

```
end
```

✓  
13/5/13

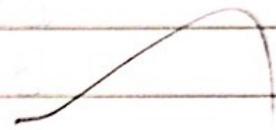
Output on the OHP:-



## \* MPT 85 Program for Triangle Wave:

ADDRESS	CODES	MNEMONICS
8000	31 C2 FF	LXI SP, FFC2
8003	3E 80	MVI A, 80H
8005	D3 2B	OUT 2BH
8007	3E 00	MVI A, 00H
8009	D3 28	OUT 28H
800B	3C	INR A
800C	FE FF	CPI FFH
800E	C2 09 80	JNZ 8009
8011	3D	DCR A
8012	D3 28	OUT 28H
8014	FE 00	CPI 00H
8016	C2 11 80	JNZ 8011
8019	C3 07 80	JMP 8007

Output on the OHP:-



# \* MPT 85 Program for Square Wave

ADDRESS	CODES	MNEMONICS
8000	31 C2 FF	LXI SP, FFC2
8003	3E 80	MVI A, 80
8005	D3 2B	OUT 2B
8007	3E 00	MVI A, 00H
8009	D3 28	OUT 28
800B	CD 18 80	CALL DELAY
800E	3E FF	MVI A, FFH
8010	D3 28	OUT 28H
8012	CD 18 80	CALL DELAY
8015	C3 07 80	JMP 8007
8018	3E FF	MVI A, FFH
801A	0D	DCR 0
801B	C2 1A 80	JNZ 801A
801E	C9	RET

After  
20/5

Output :-

CP Lab

CP Lab

CP Lab



CP Lab

CP Lab

Micro. & Interfacing LAB Notes - Akshansh

20.5.2013

# Lab - 8.9

- continued -

Q. WAP using procedure to display CP LAB at  
Top Right  
Top Left  
Center  
Bottom Right  
Bottom Left of screen.

- model small
- data  
str db 'CP Lab \$'
- code
- startup

```
mov dh, 21  
mov dl, 38  
call position  
mov dx, offset str  
call display
```

```
mov dh, 45  
mov dl, 72  
call position  
mov dx, offset str  
call display
```

```
mov dh, 0
mov dl, 0
call position
mov dx, offset str
call display
```

```
mov dh, 45
mov dl, 0
call position
mov bx dx, offset str
call display
```

```
mov dh, 0
mov dl, 72
call position
mov dx, offset str
call display
```

exit

~~position~~ position proc

```
mov ah, 2
mov bh, 0
int 10h
ret
```

position end p

```
display pro  
mov ah, 9  
int 21h  
ret  
display endp  
end z
```

✓  
~~20/5/13~~

Micro: & Interfacing Notes - Akshansh

\* Program to rotate the stepper motor in clockwise direction

Label	Address	B1	B2	B3	Mnemonic
	8000	31	C2	FF	LXI SP, FFC2H
	8003	3E	80		MVI A, 80H
	8005	D3	2B		OUT 2B.H
Loop	8007	3E	50		MVI A, 50H
	8009	D3	28		OUT 28H
	800B	CD	00	81	CALL Delay
	800E	3E	90		MVI A, 90H
	8010	D3	28		OUT 28H
	8012	CD	00	81	CALL Delay
	8015	3E	A0		MVI A, A0H
	8017	D3	28		OUT 28H
	8019	CD	00	81	CALL Delay
	801C	3E	60		MVI A, 60H
	801E	D3	28		OUT 28H
	8020	CD	00	81	CALL Delay
	8023	C3	07	80	JMP Loop

Note :- The delay routine (the speed of output) will also be required to be <sup>entered</sup> executed before executing the above program.

Output: Stepper motor rotates in clockwise direction.

\* Program to rotate the stepper motor in anti-clockwise direction .

Label	Address	B1	B2	B3	Mnemonic
	8000	31	C2		LXI SP, FFC2H
	8003	3E	80		MVI A, 80H
	8005	D3	2B		OUT 2BH
Loop	8007	3E	60		MVI A, 60H
	8009	D3	28		OUT 28H
	800B	CD	00	81	CALL Delay
	800E	3E	A0		MVI A, A0H
	8010	D3	28		OUT 28H
	8012	CD	00	81	CALL Delay
	8015	3E	90		MVI A, 90H
	8017	D3	28		OUT 28H
	8019	CD	00	81	CALL Delay
	801C	3E	50		MVI A, 50H
	801E	D3	28		OUT 28H
	8020	CD	00	81	CALL Delay
	8023	C3	07	80	JMP Loop .

The circles show the change done in previous program to reverse the direction from clockwise to anti-clockwise.

To enter value :- ~~Set~~ Set Address > 8000 >> Next >> Reset

To get o/p or execute :- Reset > Go Previous > 8000 > Exe

program  
Output : Stepper motor rotates in anti-clockwise direction .

\* Program to be written before executing the above programs :-

To set the delay of the stepper motor, i.e., the speed at which it'll rotate.

Label	Address	B1	B2	B3	Mnemonics
Delay	8100	1E	FF		MVI E, FF
	8102	CD	0A	81	CALL DLY
	8105	1D			DCR E
	8106	C2	0A	81	JNZ DLY
	8109	C9			RET
DLY	810A	0D			NOP
	810B	0D			NOP
	810C	0E	00		MVI C, 00H
Loop 2	810E	3A	50	80	LDA 8050H
	8111	3D			DCR A
	8112	C2	11	81	JNZ loop 2
	8115	0D			DCR C
	8116	C2	0E	81	JNZ loop 2
	8119	C9			RET

Note: For changing the speed of rotation, go to the address: 8050 :-

Reset > Go Prev > 8050 > ( ) > Next > Reset

↳ It can have the value from 02 (max. speed) to 20 (minimum speed).